

CSCI 210: Computer Organization

Lecture 10: Control Flow

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CS History: The If-Else Statement

- Haskell Curry and Willa Wyatt are the first people to describe performing different instructions based on the result of a previous calculation, on the Eniac in 1946
- Early assembly language instructions jumped to a new memory location based on a specific condition, were not general purpose
- Fortran (1957) specifying jumps to three locations at once, depending on whether a calculation was negative, zero, or positive, and gave it the name "if."
- Flow-matic (Grace Hopper, 1958), used comparisons between numbers and used the name "otherwise" for else
- In 1958, a German computing organization proposed an if statement that took an arbitrary Boolean statement, had an "else" case, and returned control to immediately after the if/else statement after completing the statement

Today: Program control flow

- High level languages have many ways to control the order of execution in a program: if, if-else, for loops, while loops
- Today we will look at how these higher order concepts are built out of MIPS control flow instructions

Control Flow

- Recall the basic instruction cycle
 - $IR = \text{Memory}[PC]$
 - $PC = PC + 4$
- Both branch and jump instructions change the value of the program counter

Control Flow - Instructions

- Conditional
 - **beq, bne**: compare two registers and branch depending on the comparison
 - Change the value of the program counter if a condition is true
- Unconditional
 - **j, jal, jr**: jump to a location
 - Always change the value of the program counter

Control Flow - Labels

- In assembly, we use labels to help us guide control flow. Labels can be the target of branch or jump instructions.

- Example:

```
j label
```

```
...
```

```
label:  add $t1, $t0, $t2
```

- Assemblers are responsible for translating labels into addresses.

Jump

- `j label`
 - Go directly to the label (i.e. $PC = \text{label}$)
- `jal label`
 - Go directly to the label (i.e., $PC = \text{label}$) and set the link register (we'll discuss this later)
- `jr register`
 - Go directly to the address specified in the register

High-level code

```
if (X == 0)
    X = Y + Z;
```

Assuming X, Y, and Z are integers in registers \$t0, \$t1, and \$t2, respectively, which are the equivalent assembly instructions?

A

```
        beq $t0,$zero, Label
Label:  add $t0, $t1, $t2
```

B

```
        beq $t0,$zero, Label
        add $t0, $t1, $t2
Label:
```

C

```
        bne $t0,$zero, Label
        add $t0, $t1, $t2
Label:
```

D – None of these is correct.

If ($x < y$): Set Less Than

- Set result to 1 if a condition is true
 - Otherwise, set to 0
- `slt rd, rs, rt`
 - if ($rs < rt$) $rd = 1$; else $rd = 0$;
- `slti rt, rs, constant`
 - if ($rs < \text{constant}$) $rt = 1$; else $rt = 0$;
- Use in combination with `beq`, `bne`

```
    slt $t0, $s1, $s2    # if ($s1 < $s2)
    bne $t0, $zero, L    #   branch to L
```

Branch Instruction Design

- Why not `blt`, `bge`, etc?
- Hardware for `<`, `≥`, ... slower than `=`, `≠`
 - Combining with branch involves more work per instruction
 - `beq` and `bne` are the common case

High level code often has code like this:

```
if (i < j) {  
    i = i + 1;  
}
```

Assume $\$t0$ holds i and $\$t1$ holds j . Which of the following is the correct translation of the above code to MIPS assembly (recall $\$zero$ is always 0):

<pre>slt \$t2, \$t0, \$t1 bne \$t2, \$zero, x addi \$t0, \$t0, 1 x: next instruction</pre>	<pre>slt \$t2, \$t0, \$t1 bne \$t2, \$zero, x x: addi \$t0, \$t0, 1 next instruction</pre>	<pre>slt \$t2, \$t0, \$t1 beq \$t2, \$zero, x addi \$t0, \$t0, 1 x: next instruction</pre>
--	--	--

A

B

C

D

None of the above

```
slt rd, rs, rt  
if (rs < rt) rd = 1; else rd = 0;
```

Signed vs. Unsigned

- Signed comparison: `slt`, `sltui`
- Unsigned comparison: `sltu`, `sltui`

slt vs sltu

`$s0 = 1111 1111 1111 1111 1111 1111 1111 1111`

`$s1 = 0000 0000 0000 0000 0000 0000 0000 0001`

	<code>slt \$t0, \$s0, \$s1</code>	<code>sltu \$t0, \$s0, \$s1</code>
A	<code>\$t0 = 1</code>	<code>\$t0 = 1</code>
B	<code>\$t0 = 0</code>	<code>\$t0 = 1</code>
C	<code>\$t0 = 0</code>	<code>\$t0 = 0</code>
D	<code>\$t0 = 1</code>	<code>\$t0 = 0</code>

```
slt rd, rs, rt
    if (rs < rt) rd = 1; else rd = 0;
```

Questions on BEQ, BNE, SLT?

How to access an array in a for loop

- Can't programmatically change the offset
- Need to change the *base address* instead
- Add 4 to the base address every time you want to move to the next element of the array (assuming an array of 4-byte values)

```
for (i=0; i < 10; i++) {  
    A[i] = 0;  
}
```

*Assume base address of A is in \$s0

```
        move    $t0, $zero  
        li     $t1, 40  
loop:   beq    $t0, $t1, end  
        add    $t2, $s0, $t0  
        sw    $zero, 0($t2)  
        addi   $t0, $t0, 4  
        j     loop  
end:
```


C Code

```
for (i = 0; i < 10; i++) {  
    A[i+1] = A[i];  
}
```

Assume the base address of A is in \$t0, and i is in \$t1. Each element of A is 4 bytes. What is the equivalent assembly?

```
    addi $t2, $zero, 10  
    add $t1, $zero, $zero  
for: bne $t1, $t2, end  
    lw $t3, $t1($t0)  
    addi $t1, $t1, 1  
    sw $t3, $t1($t0)  
    j for  
end:
```

A

```
    addi $t2, $zero, 40  
    add $t1, $zero, $zero  
for: beq $t1, $t2, end  
    add $t4, $t0, $t1  
    lw $t3, 0($t4)  
    addi $t1, $t1, 4  
    add $t4, $t0, $t1  
    sw $t3, 0($t4)  
    j for  
end:
```

B

```
    addi $t2, $zero, 10  
    add $t1, $zero, $zero  
    bne $t1, $t2, end  
    add $t4, $t0, $t1  
    lw $t3, 0($t4)  
    addi $t1, $t1, 1  
    add $t4, $t0, $t1  
    sw $t3, 0($t4)  
end:
```

C

D – More than one of these

E – None of these

C Code

```
if (X == 0)
    X = Y + Z;
else
    X = Z + Z;
```

Assuming X, Y, and Z are integers in registers \$t0, \$t1, and \$t2, respectively, which are the equivalent assembly instructions?

```
bne $t0, $zero, false
add $t0, $t1, $t2
false: add $t0, $t2, $t2
```

A

```
bne $t0, $zero, false
add $t0, $t1, $t2
j endif
false: add $t0, $t2, $t2
endif:
```

B

```
bne $t0, $zero, false
j endif
add $t0, $t1, $t2
false: add $t0, $t2, $t2
endif:
```

C

D – None of the above

C Code

```
while (i < 10) {  
    i = i + 1;  
}
```

Assume i is in \$t0. What is the equivalent assembly?

```
slti rt, rs, imm  
if (rs < imm) rd = 1; else rt = 0;
```

```
w: slti $t2, $t0, 10  
    beq $t2, $zero, end  
    addi $t0, $t0, 1  
    j w  
end:
```

A

```
w: slti $t2, $t0, 10  
    beq $t2, $zero, end  
    addi $t0, $t0, 1  
end:
```

B

```
slti $t2, $t0, 10  
w: beq $t2, $zero, end  
    addi $t0, $t0, 1  
    j w  
end:
```

C

D – More than one of these

E – None of these

Reading

- Next lecture: Procedures
 - Section 2.9
- Problem set: Due Friday
- Lab 2: Due Monday